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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/085,455	02/27/2002	Motohiro Kawahito	JP920000420US1	1801
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IBM CORPORATION			PHAM, CHRYSTINE	
INTELLECTUAL PROPERTY LAW DEPT. P.O. BOX 218 - 39-254 YORKTOWN HEIGHTS, NY 10598			ART UNIT	PAPER NUMBER
			2122	
			DATE MAILED: 01/26/2005	

Please find below and/or attached an Office communication concerning this application or proceeding.

		Application N .	Applicant(s)			
Office Action Summary		10/085,455	KAWAHITO ET AL.			
		Examiner	Art Unit			
		Chrystine Pham	2122			
The MAILING DATE of this communication appears on the c ver sheet with the correspondence address Period for Reply						
THE - External after - If the - If NC - Failu Any	ORTENED STATUTORY PERIOD FOR REPL MAILING DATE OF THIS COMMUNICATION. nsions of time may be available under the provisions of 37 CFR 1. SIX (6) MONTHS from the mailing date of this communication. e period for reply specified above is less than thirty (30) days, a repl period for reply is specified above, the maximum statutory period re to reply within the set or extended period for reply will, by statut reply received by the Office later than three months after the mailing datent term adjustment. See 37 CFR 1.704(b).	136(a). In no event, however, may a reply be tin ply within the statutory minimum of thirty (30) day I will apply and will expire SIX (6) MONTHS from te, cause the application to become ABANDONE	nely filed s will be considered timely. the mailing date of this communication. D (35 U.S.C. § 133).			
Status						
1)⊠	1) Responsive to communication(s) filed on <u>27 February 2002</u> .					
2a) <u></u>	This action is FINAL . 2b)⊠ Thi	is action is non-final.				
3)□	Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under <i>Ex parte Quayle</i> , 1935 C.D. 11, 453 O.G. 213.					
Dispositi	ion of Claims					
5)□ 6)⊠ 7)⊠	4) ☐ Claim(s) 1-16 is/are pending in the application. 4a) Of the above claim(s) is/are withdrawn from consideration. 5) ☐ Claim(s) is/are allowed. 6) ☐ Claim(s) 1-16 is/are rejected. 7) ☐ Claim(s) 2 is/are objected to. 8) ☐ Claim(s) are subject to restriction and/or election requirement.					
Applicati	ion Papers					
10)⊠	The specification is objected to by the Examin The drawing(s) filed on 27 February 2002 is/a Applicant may not request that any objection to the Replacement drawing sheet(s) including the correct The oath or declaration is objected to by the Examin Theorem 1.	re: a)⊠ accepted or b)□ objecte e drawing(s) be held in abeyance. Sec ction is required if the drawing(s) is ob	e 37 CFR 1.85(a). jected to. See 37 CFR 1.121(d).			
Priority (under 35 U.S.C. § 119					
12) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f). a) All b) Some * c) None of: 1. Certified copies of the priority documents have been received. 2. Certified copies of the priority documents have been received in Application No 3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)). * See the attached detailed Office action for a list of the certified copies not received.						
Attachmen	nt(s)					
2) Notice 3) Information	ce of References Cited (PTO-892) ce of Draftsperson's Patent Drawing Review (PTO-948) mation Disclosure Statement(s) (PTO-1449 or PTO/SB/08 er No(s)/Mail Date <u>17 December 2004</u> .	4) Interview Summary Paper No(s)/Mail D 5) Notice of Informal F 6) Other:				

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DETAILED ACTION

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This action is responsive to application 10/085455 filed on February 27th 2002. Claims 1-16 are presented for examination. Priority date February 28th 2001 has been considered.

Specification

2. The lengthy specification has not been checked to the extent necessary to determine the presence of all possible minor errors. Applicant's cooperation is requested in correcting any errors of which applicant may become aware in the specification.

Claim Objections

3. Claim 2 is objected to because of the following informalities: It contains an improper reference to a base claim (see "according to claim 2", line 1). Appropriate correction is required. For compact prosecution of the claims, limitation "according to claim 2" has been interpreted as "according to claim 1".

Claim Rejections - 35 USC § 102

4. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless -

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

5. Claims 1-16 are rejected under 35 U.S.C. 102(b) as being anticipated by Faiman, Jr. (USP 5836014), hereinafter, *Faiman*.

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Faiman teaches a program optimization method (e.g., see *optimizing 26* FIG.1 & associated text) for translating, into machine code (e.g., see *object code 23*, executable image, target computer 25 FIG.1 & associated text; see *object code images, machine language* col.2:49-55), source code for a program written in a programming language (e.g., see *source code 21* FIG.1 & associated text; see *compiler front end* col.2:5-16), and for optimizing said program (e.g., see *optimizing 26* FIG.1 & associated text) comprising the steps of:

- o performing an analysis to determine whether the execution speed of said program can be increased (e.g., see *optimize speed of execution* col.2:28-35) by fixing, in a specific state (e.g., see *K-folding, KFOD routine, constant* col.4:14-21), a parameter for a predetermined command in said program (e.g., see *expressions* col.4:14-21; col.22:29-34); and
- o employing results of said analysis for the generation, in said program, of a path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35).

Claim 2

The rejection of base claim 1 is incorporated. *Faiman* further teaches wherein said step of generating a path includes the steps of:

- executing said program and obtaining statistical data for the appearance frequency of each
 available state (e.g., see analyzing induction variables col.3:65-col.4:5; see detection of induction
 variables col.18:40-67) wherein, according to said results of said analysis, said parameter of said
 predetermined command may be set (e.g., see inductive expressions, multiplications, additions
 col.4:1-10); and
- employing said obtained statistical data (e.g., see inductive expressions col.18:64-col.19:50) to generate said path.

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Claim 3

Faiman teaches a program optimization method (e.g., see *optimizing 26* FIG.1 & associated text) for translating, into machine code (e.g., see *object code 23*, executable image, target computer 25 FIG.1 & associated text; see *object code images, machine language* col.2:49-55), the source code for a program written in a programming language (e.g., see *source code 21* FIG.1 & associated text; see *compiler front end* col.2:5-16), and for optimizing said program (e.g., see *optimizing 26* FIG.1 & associated text) comprising the steps of:

- o executing a program to obtain statistical data for an appearance frequency of each available state (e.g., see *analyzing induction variables* col.3:65-col.4:5; see *detection of induction variables* col.18:40-67) in which a parameter of a predetermined command in said program may be set (e.g., see *inductive expressions, multiplications, additions* col.4:1-10); and
- o employing said obtained statistical data (e.g., see *inductive expressions* col.18:64-col.19:50) to generate a machine language program that includes, as the compiling results, a path (e.g., see *constant expression evaluation routine, runtime, object code image, Kfold routine* col.22:6-20; see *generating code* col.22:67-col.23:8; see *machine code, constant expression evaluation routine* col.23:32-35) along which said parameter of said predetermined command (e.g., see *expressions* col.4:14-21; col.22:29-34) is fixed in a specific state (e.g., see *K-folding, KFOD routine, constant* col.4:14-21).

Claim 4

The rejection of base claim 3 is incorporated. *Faiman* further teaches comprising a step of: generating a machine language program that does not include, as a compiling result, a path along which said parameter of said predetermined command is fixed in a specific state (e.g., see *conditional stores* col.20:8-14).

Faiman teaches a program optimization method (e.g., see *optimizing 26* FIG.1 & associated text) for translating, into machine code (e.g., see *object code 23*, executable image, target computer 25 FIG.1 & associated text; see *object code images, machine language* col.2:49-55), the source code for a program written in an object-oriented programming language (e.g., see *source code 21*, *compiler front end 20* FIG.1 & associated text; see *front end*, C++ col.2:65-col.3:10; see *compiler front end* col.2:5-16), and for optimizing said program (e.g., see *optimizing 26* FIG.1 & associated text) comprising the steps of:

- o detecting one command, of the commands in said program, for which a method call destination can be identified (e.g., see *tuples*, *routine call*, *procedure calls* col.14:47-65), and for which the processing speed can be increased by identifying said method call destination (e.g., see *tuples* col.12:59-65; see *flow graphs*, *procedure calls*, *variable*, *memory locations* col.13:43-55; ; see *optimize speed of execution*, *flow graph* col.2:28-34); and
- o generating a path (e.g., see constant expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35) wherefor said method call destination for said detected command is limited in order to increase the processing speed of said command (e.g., see Restrictions, tuple col.55:53-66).

Claim 6

Faiman teaches a program optimization method (e.g., see *optimizing 26* FIG.1 & associated text) for translating, into machine code (e.g., see *object code 23, executable image, target computer 25* FIG.1 & associated text; see *object code images, machine language* col.2:49-55), the source code for a program written in a programming language (e.g., see *source code 21* FIG.1 & associated text; see *compiler front end* col.2:5-16), and for optimizing said program (e.g., see *optimizing 26* FIG.1 & associated text) comprising the steps of:

o detecting one command, of the commands in said program (e.g., see *expressions* col.4:14-21; col.22:29-34), for which a variable can be limited to a predetermined constant value (e.g., see *K*-

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folding, KFOD routine, constant col.4:14-21), and for which the processing speed can be increased by limiting said variable to said constant value (e.g., see *optimize speed of execution* col.2:28-35); and

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o generating a path along which said constant value of said variable of said detected command is fixed (e.g., see constant expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35).

Claim 7

Faiman teaches a compiler (e.g., see *compiler front end 20* FIG.1 & associated text) for translating into machine code (e.g., see *object code 23, executable image, target computer 25* FIG.1 & associated text; see *object code images, machine language* col.2:49-55) the source code for a program written in a programming language (e.g., see *source code 21* FIG.1 & associated text; see *compiler front end* col.2:5-16), and for optimizing the resultant program (e.g., see *optimizing 26* FIG.1 & associated text) comprising:

- o an impact analysis unit for performing an analysis to determine how much (e.g., see effects 42, dependencies 43 FIG.4 & associated text) the execution speed of said program can be increased (e.g., see optimize speed of execution col.2:28-35) by fixing, in a specific state (e.g., see K-folding, KFOD routine, constant col.4:14-21), a parameter of a predetermined command in said program (e.g., see expressions col.4:14-21; col.22:29-34); and
- o a specialization unit for employing the analysis results obtained by said impact analysis unit to generate, in said program, a specialized path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant expression evaluation routine*, runtime, object code image, Kfold routine col.22:6-20; see *generating code* col.22:67-col.23:8; see *machine code*, *constant expression evaluation routine* col.23:32-35).

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The rejection of base claim 7 is incorporated. Faiman further teaches:

o a data specialization selector for, when said program is executed, obtaining statistical data for the appearance frequency of each state obtained by said impact analysis unit (e.g., see *analyzing induction variables* col.3:65-col.4:5; see *detection of induction variables* col.18:40-67), and for determining the state in which said parameter of said predetermined command is to be set (e.g., see *inductive expressions, multiplications, additions* col.4:1-10),

o wherein said specialization unit generates a specialized path along which said parameter of said predetermined command is fixed in a state determined by said data specialization selector (e.g., see *inductive expressions* col.18:64-col.19:50).

Claim 9

The rejection of base claim 8 is incorporated. *Faiman* further teaches wherein, in accordance with the state of said program at execution, said specialization unit generates, in said program, a branching process for selectively performing a specialized path and an unspecialized path (e.g., see *conditional branch*, cases col.23:20-40); and wherein, while taking into account a delay due to the insertion of said branching process (e.g., see *Delayed Actions*, passes col.25:30-40), said data specialization selector determines a state in which said parameter of said predetermined command is fixed (e.g., see *inductive expressions* col.18:64-col.19:50).

Claim 10

Faiman teaches a computer (e.g., see FIGS.1, 2 & associated text) comprising:

- o an input device for receiving source code for a program (e.g., see source code 21, compiler front end 20, shell 11 FIG.1 & associated text; col.6:14-21);
- o a compiler (e.g., compiler front end 20 FIG.1 & associated text) for translating said source code to compile said program (e.g., see compiler front end col.2:5-16) and for converting said compiled program into machine language code (e.g., see object code 23, executable image FIG.1 & associated text; see object code images, machine language col.2:49-55); and

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o a processor for executing said machine language code (e.g., see CPU 14 FIG.2 & associated text; see target computer 25 FIG.1 & associated text),

- o wherein said compiler includes
- o means for performing an analysis to determine whether the execution speed of said program can be improved (e.g., see *optimize speed of execution* col.2:28-35) by fixing in a specific state (e.g., see *K-folding, KFOD routine, constant* col.4:14-21) a parameter of a predetermined command in said program (e.g., see *expressions* col.4:14-21; col.22:29-34), and
- o means for generating in said program, based on the analysis results, a path along which said parameter of said predetermined command is fixed in said specific state and for compiling said program (e.g., see *constant expression evaluation routine*, *runtime*, *object code image*, *Kfold routine* col.22:6-20; see *generating code* col.22:67-col.23:8; see *machine code*, *constant expression evaluation routine* col.23:32-35), and
- o wherein said compiler outputs, as the compiled results, said machine language code that includes said path along which the state of said parameter is fixed (e.g., see *constant expression* evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35).

Claim 11

Faiman teaches a computer (e.g., see FIGS.1, 2 & associated text) comprising:

- o an input device, for receiving source code for a program (e.g., see source code 21, compiler front end 20, shell 11 FIG.1 & associated text; col.6:14-21);
- o a compiler (e.g., compiler front end 20 FIG.1 & associated text), for translating said source code to compile said program (e.g., see compiler front end col.2:5-16) and for converting said compiled program into machine language code (e.g., see object code 23, executable image FIG.1 & associated text; see object code images, machine language col.2:49-55); and
- a processor, for executing said machine language code (e.g., see CPU 14 FIG.2 & associated text; see target computer 25 FIG.1 & associated text),

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o wherein said compiler includes

o means for obtaining statistical data for the appearance frequency of each available state wherein a parameter for a predetermined command in said program may be set when said program is executed (e.g., see *analyzing induction variables* col.3:65-col.4:5; see *detection of induction variables* col.18:40-67), and for employing said statistical data to determine a state in which said parameter of said predetermined command is to be fixed (e.g., see *inductive expressions*, *multiplications*, *additions* col.4:1-10), and

- o means for generating a specialized path along which said parameter of said predetermined command is fixed in said determined state, and for compiling said program (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35), and
- o wherein said compiler outputs, as the compiled results, said program as said machine language code that includes said specialized path (e.g., see *constant expression evaluation routine*, runtime, object code image, Kfold routine col.22:6-20; see *generating code* col.22:67-col.23:8; see *machine code*, *constant expression evaluation routine* col.23:32-35).

Claim 12

The rejection of base claim 11 is incorporated. *Faiman* further teaches comprising: said compiler further includes means for compiling said program without generating a specialized path, wherein, when said state of said parameter to be fixed can not be determined, said means for determining the state of said parameter of said predetermined command outputs, as compiled results, said program in said machine language code, which is generated by said means for compiling said program without generating said specialized path, that does not include said specialized path (e.g., see *conditional stores* col.20:8-14).

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Faiman teaches a support program, for controlling a computer to support generation of a program (e.g., see machine code, constant expression evaluation routine col.23:32-35), which permits said computer to perform:

- o a function for performing an analysis to determine whether the execution speed of said program can be increased (e.g., see *optimize speed of execution* col.2:28-35) by fixing a parameter of a predetermined command (e.g., see *expressions* col.4:14-21; col.22:29-34) of said computer program in a specific state (e.g., see *K-folding, KFOD routine, constant* col.4:14-21); and
- o a function for generating in said program, based on the analysis results, a path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35).

Claim 14

Faiman teaches a support program, for controlling a computer to support generation of a program (e.g., see machine code, constant expression evaluation routine col.23:32-35), which permits said computer to perform:

- o a function for executing said program and obtaining statistical data for the appearance frequency of each available state (e.g., see analyzing induction variables col.3:65-col.4:5; see detection of induction variables col.18:40-67) wherein said parameter of said predetermined command of said program may be set (e.g., see inductive expressions, multiplications, additions col.4:1-10); and
- o a function for generating in said program, based on said statistical data, a path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35; inductive expressions col.18:64-col.19:50).

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Claim 15

Faiman teaches a storage medium (e.g., see memory 15, 17 FIG.2 & associated text) on which input means of a computer stores a computer-readable support program (e.g., see machine code, constant expression evaluation routine col.23:32-35), for controlling said computer to support generation of a program, that permits said computer to perform:

- o a function for performing an analysis to determine whether the execution speed of said program can be increased (e.g., see *optimize speed of execution* col.2:28-35) by fixing a parameter of a predetermined command (e.g., see *expressions* col.4:14-21; col.22:29-34) of said computer program in a specific state (e.g., see *K-folding, KFOD routine, constant* col.4:14-21); and
- o a function for generating in said program, based on the analysis results, a path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35).

Claim 16

Faiman teaches a storage medium (e.g., see memory 15, 17 FIG.2 & associated text) on which input means of a computer stores a computer-readable support program (e.g., see machine code, constant expression evaluation routine col.23:32-35), for controlling said computer to support generation of a program, that permits said computer to perform:

o a function for executing said program and obtaining statistical data for the appearance frequency of each available state (e.g., see *analyzing induction variables* col.3:65-col.4:5; see *detection of induction variables* col.18:40-67) wherein said parameter of said predetermined command of said program may be set (e.g., see *inductive expressions, multiplications, additions* col.4:1-10); and

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o a function for generating in said program, based on said statistical data, a path along which said parameter of said predetermined command is fixed in said specific state (e.g., see *constant* expression evaluation routine, runtime, object code image, Kfold routine col.22:6-20; see generating code col.22:67-col.23:8; see machine code, constant expression evaluation routine col.23:32-35; inductive expressions col.18:64-col.19:50).

Conclusion

- 6. The prior art made of record and not relied upon is considered pertinent to applicant's disclosure.
- 7. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Chrystine Pham whose telephone number is 571-272-3702. The examiner can normally be reached on Mon-Fri, 8:30am-5pm.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Tuan Q Dam can be reached on 571-272-3695. The fax phone number for the organization where this application or proceeding is assigned is 703-872-9306.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

January 21, 2005

SUPERVISORY PATENT EXAMINER